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# The Dynamic Behavior of a Data Dissemination Protocol for Network Programming at Scale

*Jonathan W. Hui, David Culler*

Presented by:

Thomas Repantis

trep@cs.ucr.edu

# Overview

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Deluge, a reliable data dissemination protocol for propagating large data objects from one or more source nodes to many other nodes over a multihop, wireless sensor network.

1. Introduction
2. Protocol Description
3. Experimental Evaluation
4. Conclusions
5. Comparison with P2P

# Network Programming

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- Deluge's goal is the propagation of complete binary images to the nodes of the network.
- Network programming (the programming of nodes by disseminating code over the network) is useful:
  - Application requirements can change.
  - Thousands of nodes embedded.
  - Also useful for debugging and testing.

# Requirements

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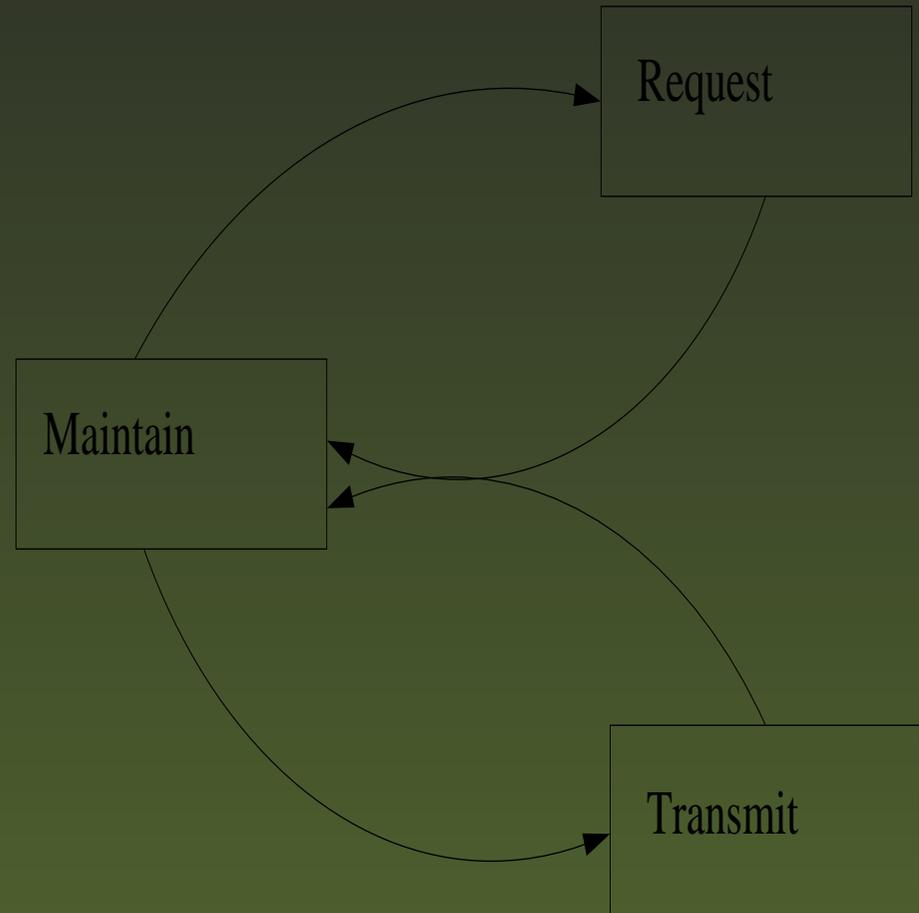
Issues in reliable dissemination of large data objects from few to many nodes of a multi-hop WSN:

- Large data objects vs constrained storage hierarchy:  
Packet (32 bytes) « RAM (4K) « program (128K) < external flash (512K)
- High node density (hinders transmission).
- Complete reliability required, even with highly lossy links.
- All nodes need to receive the newest code version.
- Fast propagation is desired.

# Deluge Protocol Overview

Nodes follow strictly local rules:

- Advertise
- Request data
- Broadcast



# Epidemic Protocol

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- Each node occasionally advertises (by broadcasting) the most recent version of the data object it has available.
- If a node receives an advertisement from an older node, it responds with its object profile.
- The older node determines from the object profile which portions of its data need updating and requests them from any neighbor that advertises the availability of the needed data.
- Nodes receiving the request then broadcast any requested data.
- Nodes then advertise newly received data in order to propagate it further.

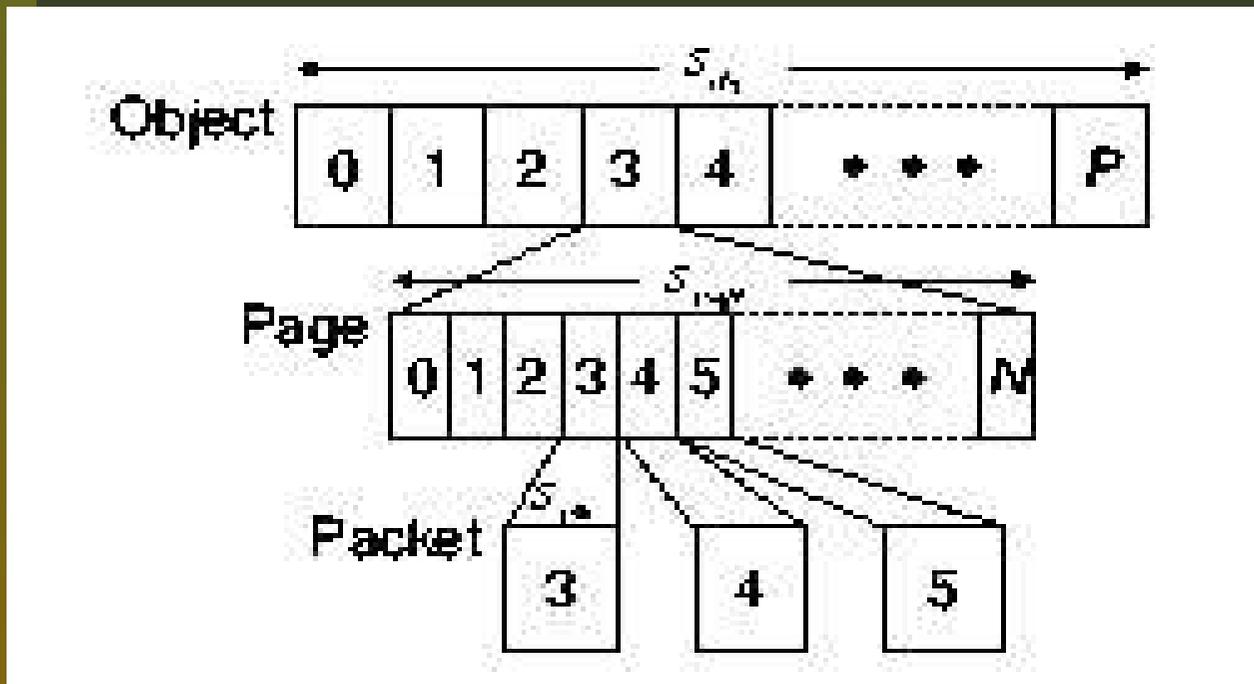
# Protocol Properties

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- Density-aware (redundant advertisements and requests are suppressed).
- Robust to asymmetric links (three-phase handshaking).
- Switches between quick propagation and quiet phases (dynamically adjustable advertisement rate).
- Minimizes the set of nodes concurrently broadcasting data within a cell.
- Allows parallel transfers of data (spatial multiplexing).

# Data Representation 1/2

CRCs at page and packet level ensure correct transmission.



# Data Representation 2/2

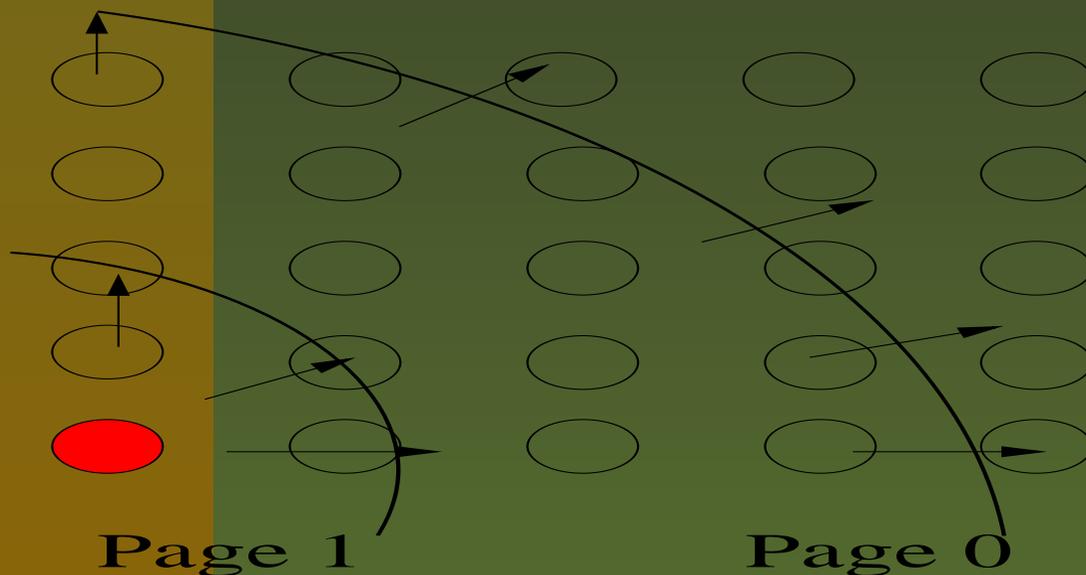
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Advantages of fixed-size pages:

- Reduced memory requirements for maintaining state about which packets are needed.
- Enables efficient **incremental updates** using:
  - A monotonically increasing **version number**  $v$  for every update.
  - An **age vector**  $a$  describes how old a page is.
  - An **object profile** is defined by the tuple  $(v,a)$ .
- Allows for **spatial multiplexing**.

# Spatial Multiplexing

- Propagation in “waves” (pipelining the page transfers).
- Exploits the limited range of radio to allow for concurrent broadcasts.
- Reduces the completion time from  $O(d \cdot S_{obj})$  to  $O(d + S_{obj})$ .



# Maintain State

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- Advertise version and fraction of image that is complete (version number and highest page available, since nodes request pages sequentially).
- Deluge uses Trickle to control the advertisement broadcasts.
- Uses suppression to decrease the advertisement rate as neighbors increase.
- Bounds advertisement rate independent of node density (by suppressing the transmission of redundant advertisements).
- Allows for quick propagation during an upgrade and low resource consumption in the steady state by decreasing and increasing the advertisement rate respectively.

# Maintain State Transitions

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- Transition to **Transmit**:  
When receiving a request.
- Transition to **Request**:  
When receiving an advertisement with newer data,  
unless a request or data packet was recently  
overheard.

# Request State

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- Transmit a request after random backoff time.
- Suppress if any similar requests of data packets are overheard during the backoff period.
- Minimize senders by unicasting requests to the node that advertised.
- Transition to maintain after receiving all packets of a page or after  $k$  requests (to protect against asymmetric lossy links).

# Transmit State

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- Transmit all requested packets.
- Packets are sent in round-robin order to provide fairness among requesters.
- Transition to maintain after all requested packets have been transmitted and no new packets have been requested.

# Experimental Evaluation

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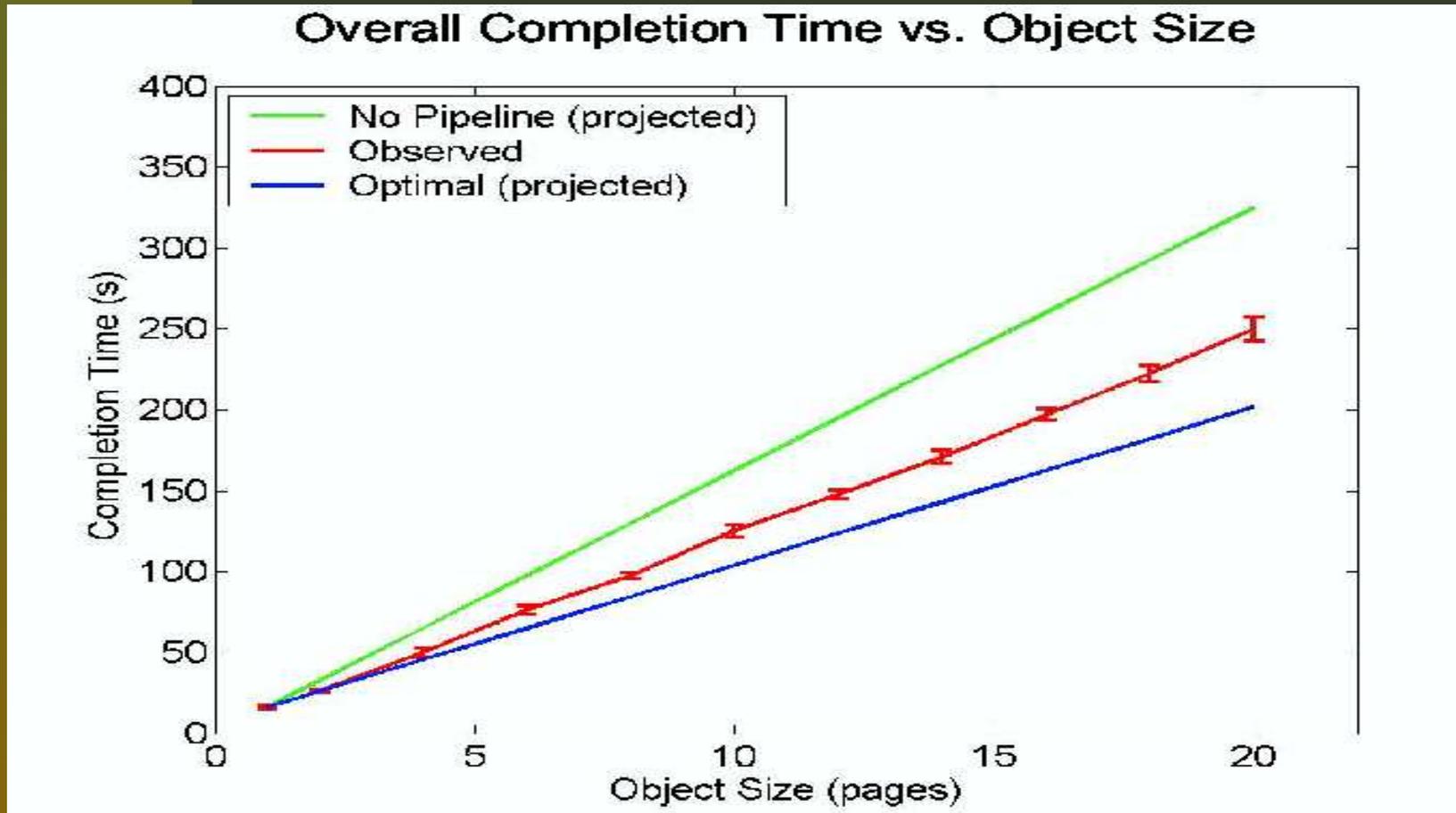
- Real-world deployment (77 Mica2-dot nodes, in a building).
- Simulation (TOSSIM, up to 800 nodes, in a square grid).

## Metrics:

- Complete Reliability
- Completion Time
- RAM Usage
- Energy Consumption

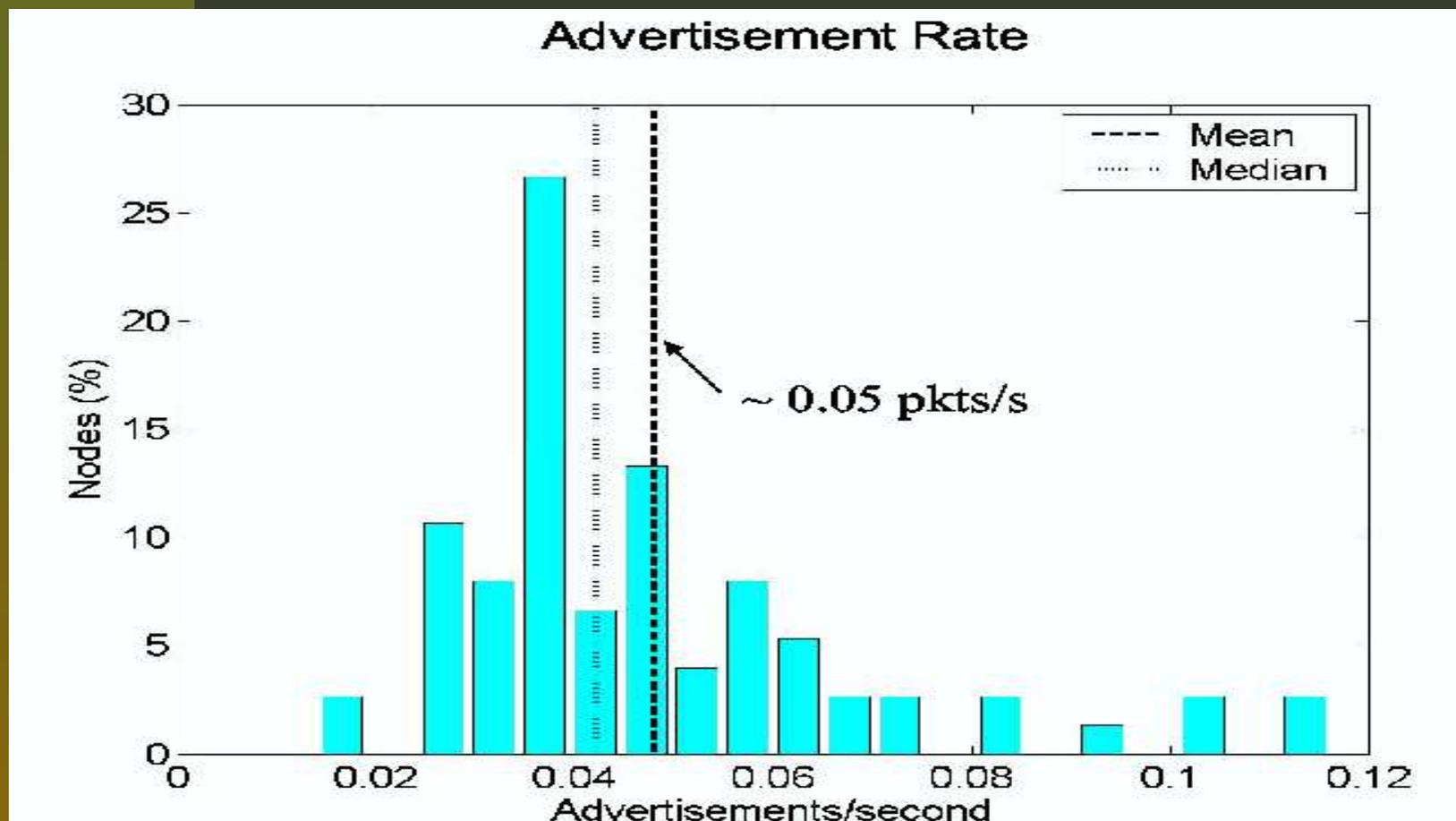
# Real World Completion Time

Completion time is linear with the object size.



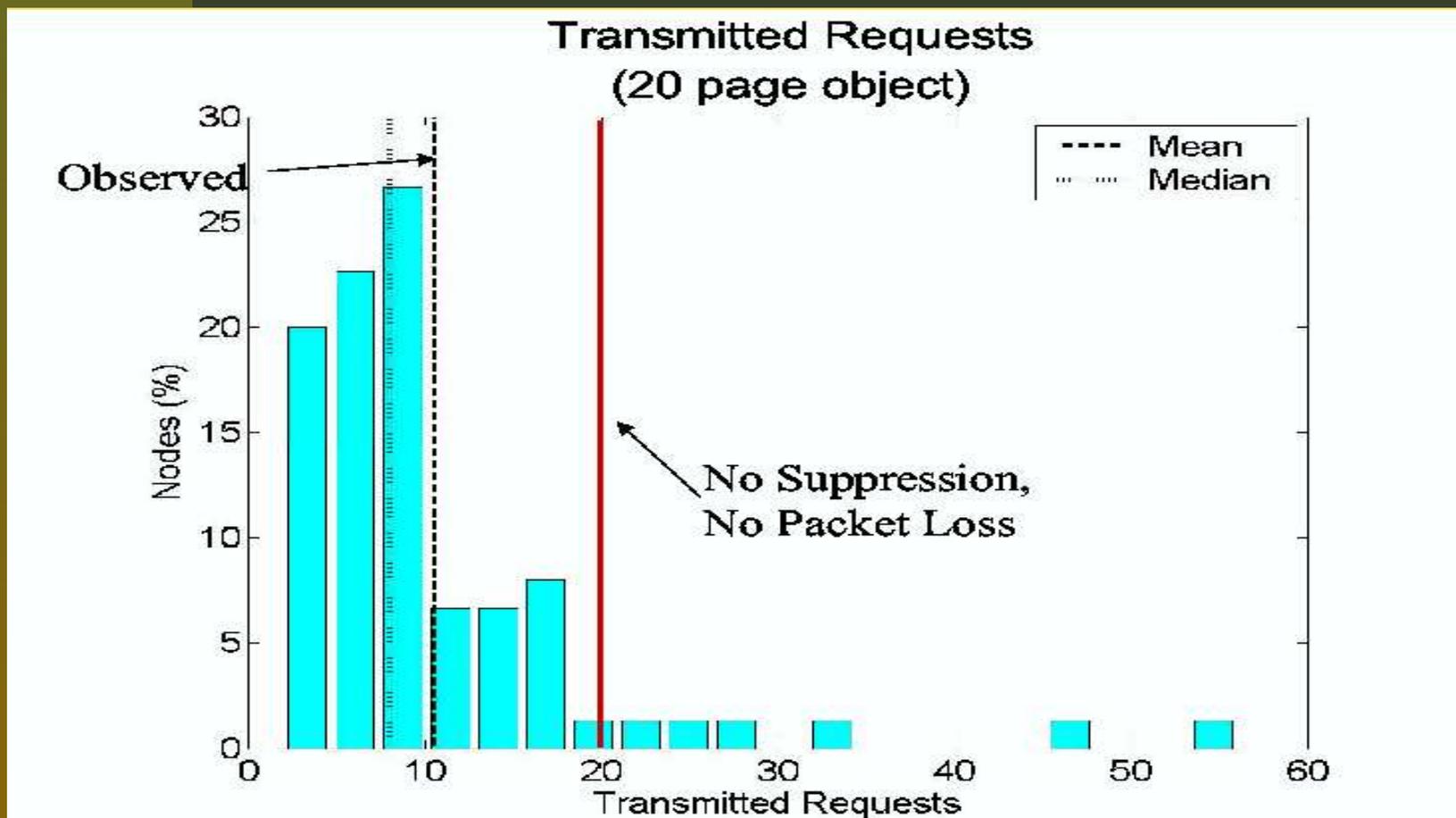
# Real World Advertisement Rate

Without suppression it would have been ten times more.



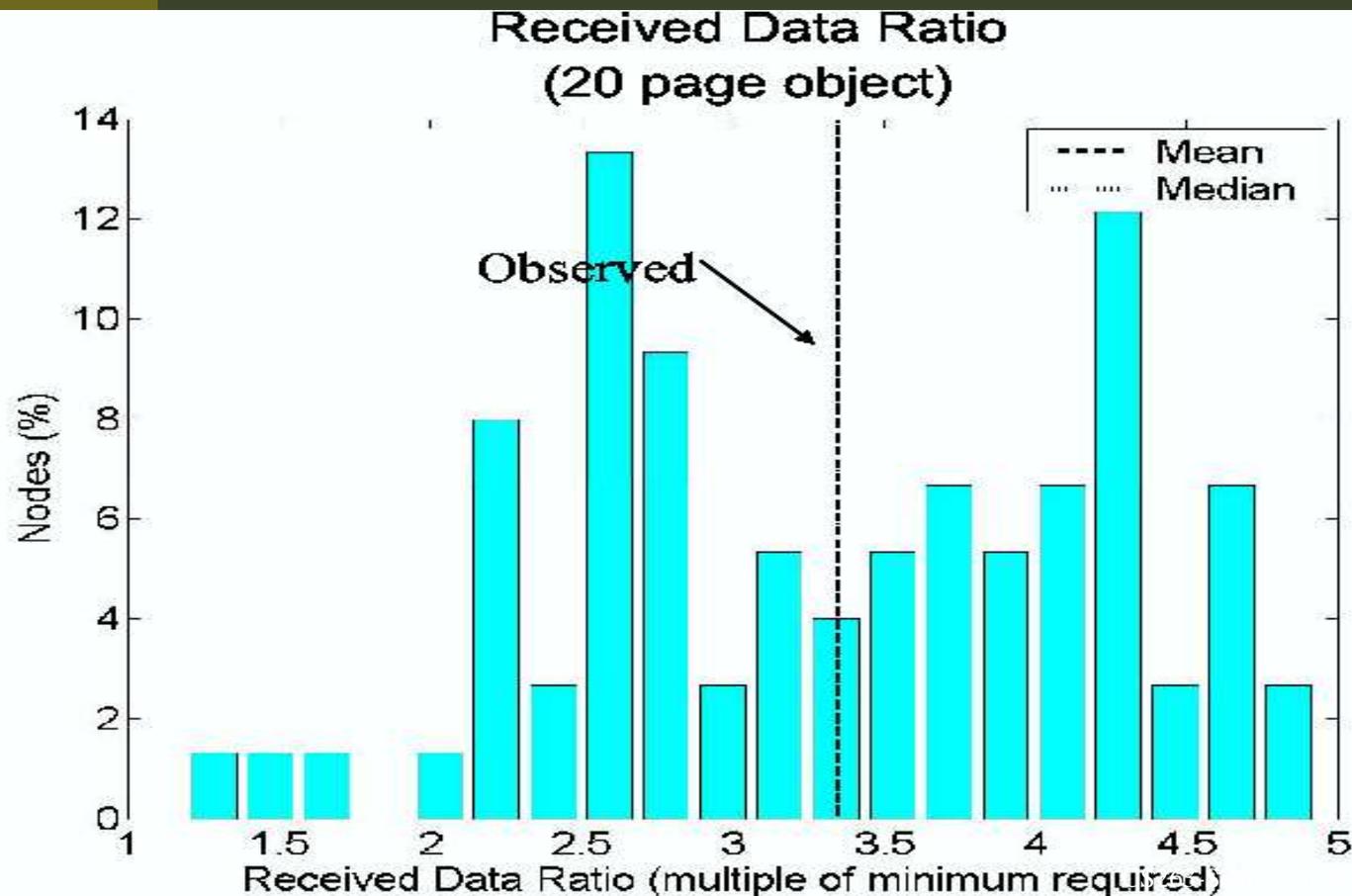
# Real World Transmitted Requests

Many nodes took advantage of the broadcast channel and did not need to transmit requests to receive entire pages.



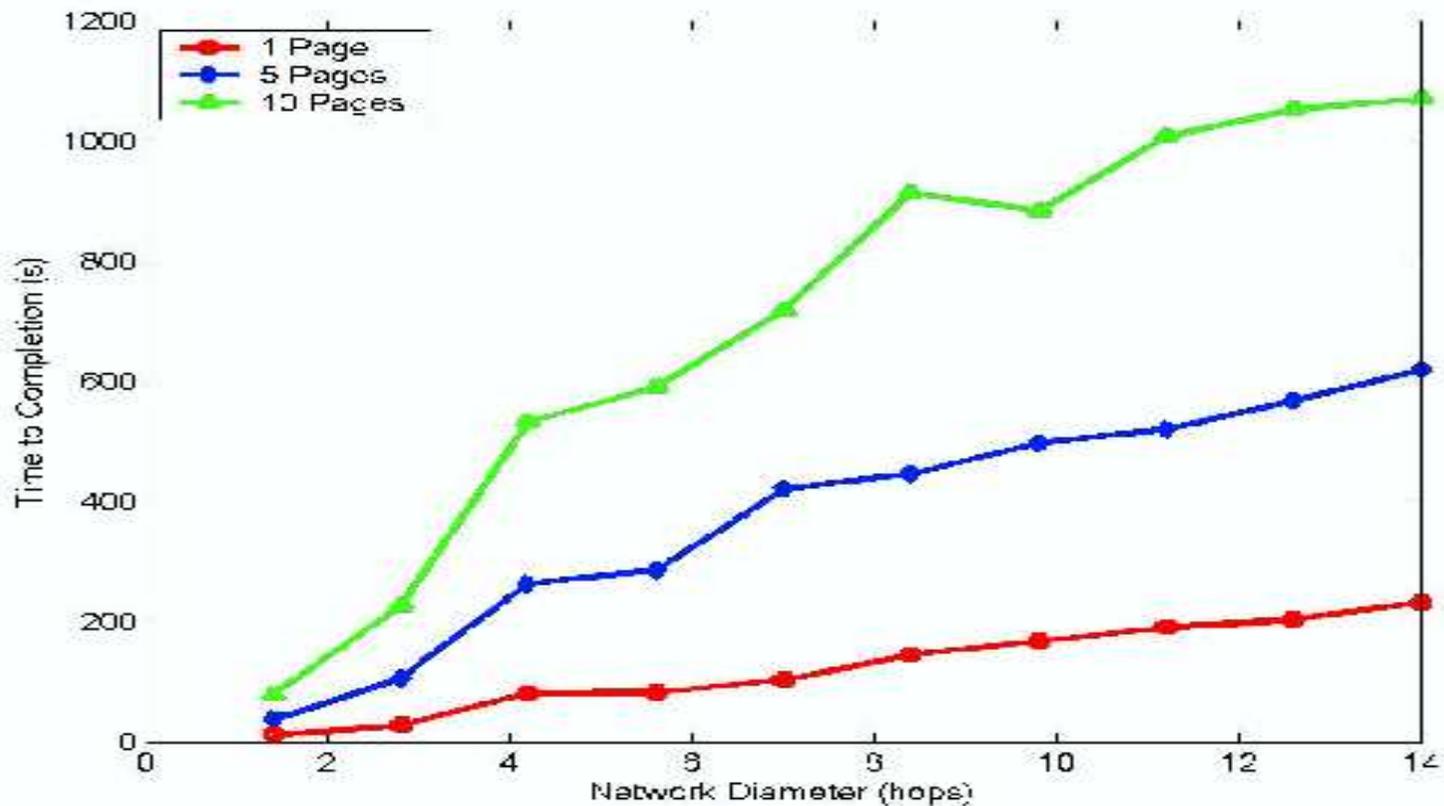
# Real World Data Redundancy

On average a node receives about 3.35 times the minimum number of required data packets, due to the single-channel, broadcast network.



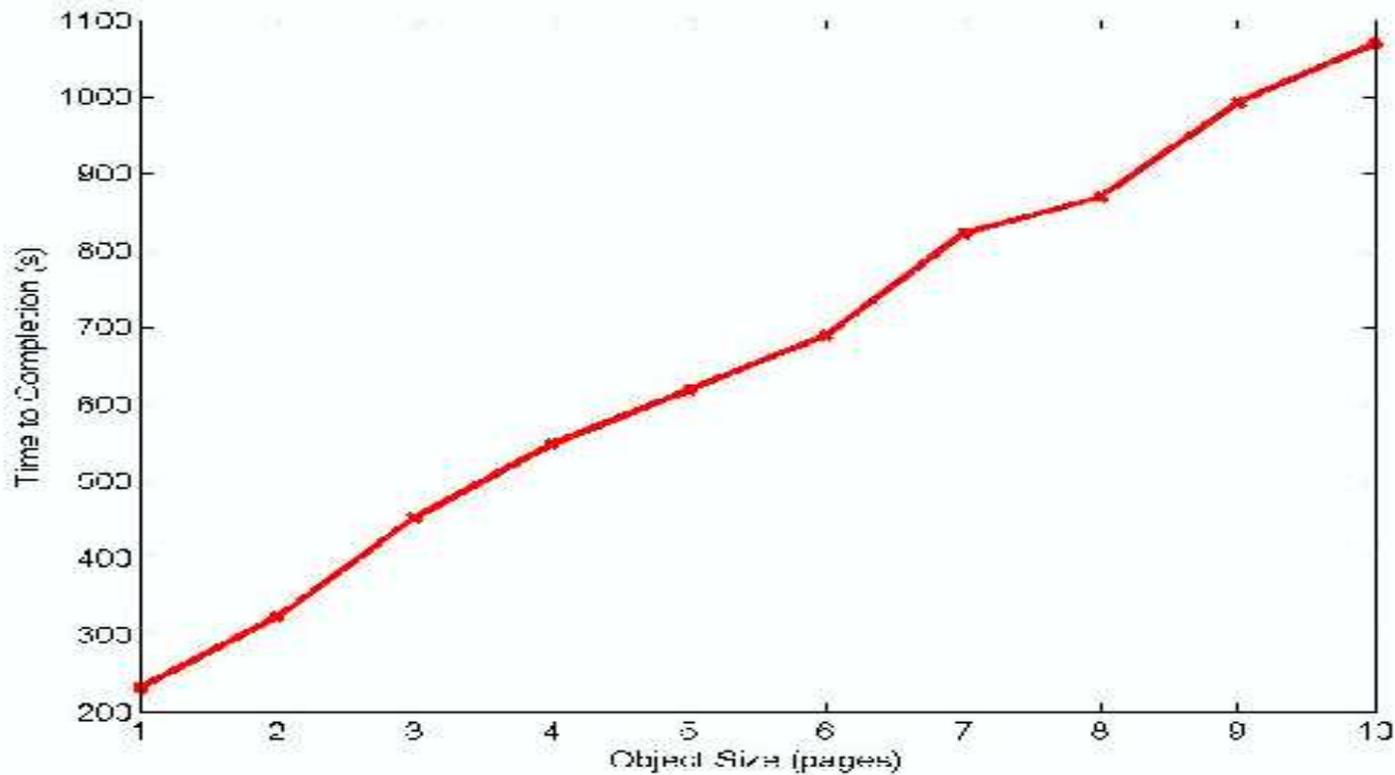
# Simulated Completion Time for Varying Network Diameter

Spatial multiplexing is effective (propagation time is the sum of the network diameter and object size).



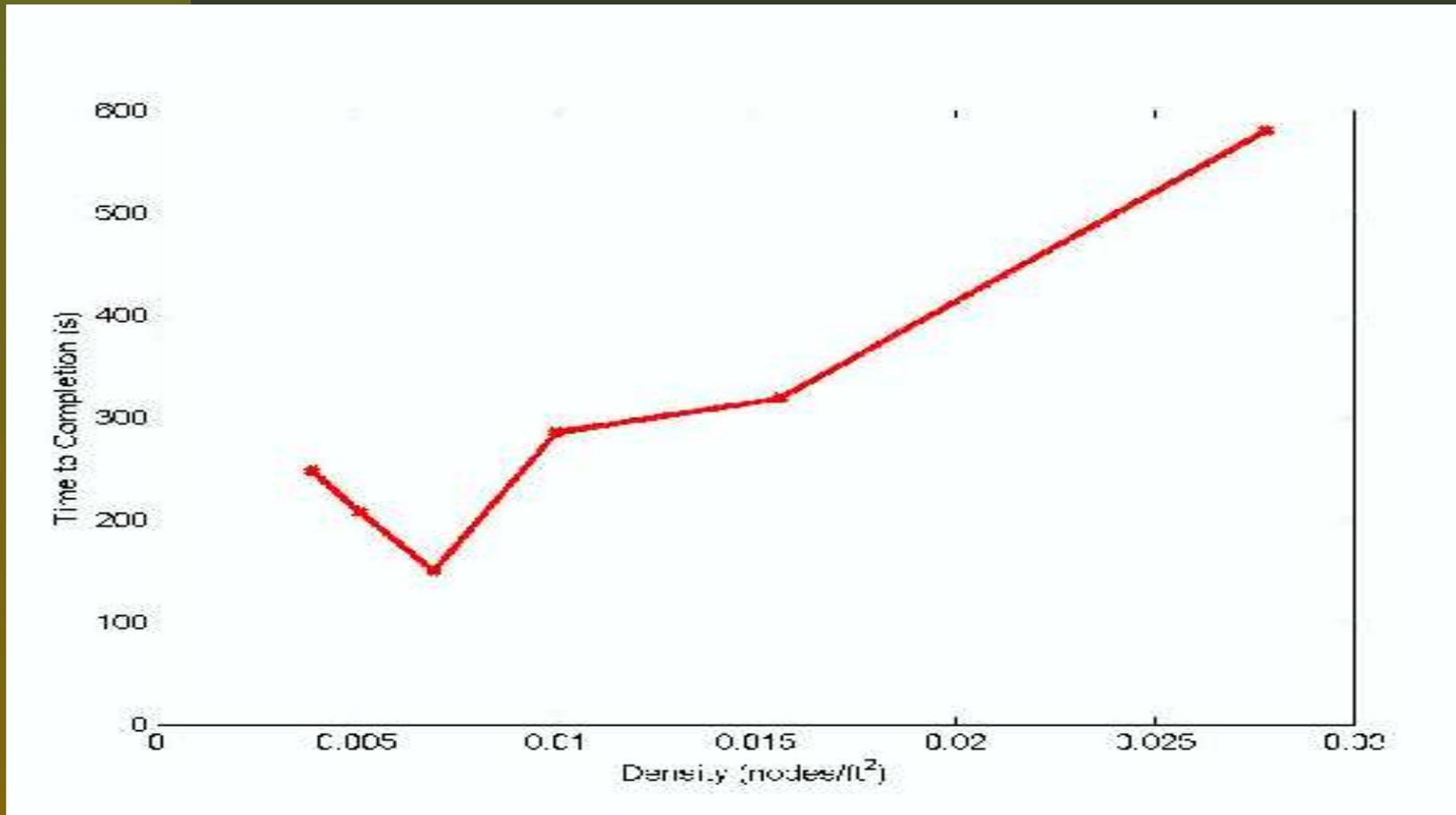
# Simulated Completion Time for Varying Object Size

Completion time is linear with the object size.



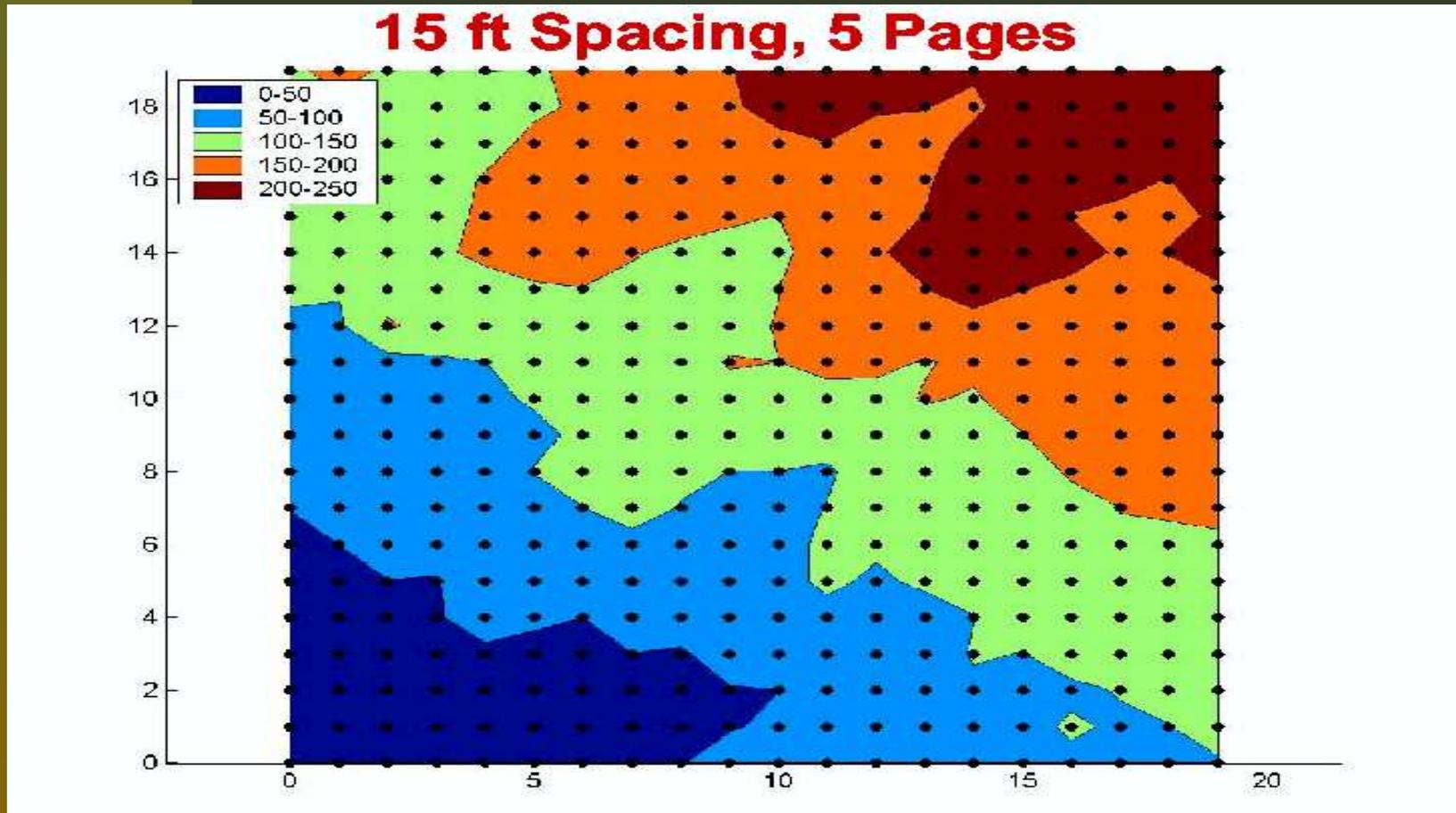
# Simulated Completion Time for Varying Density

Completion time increases with density.



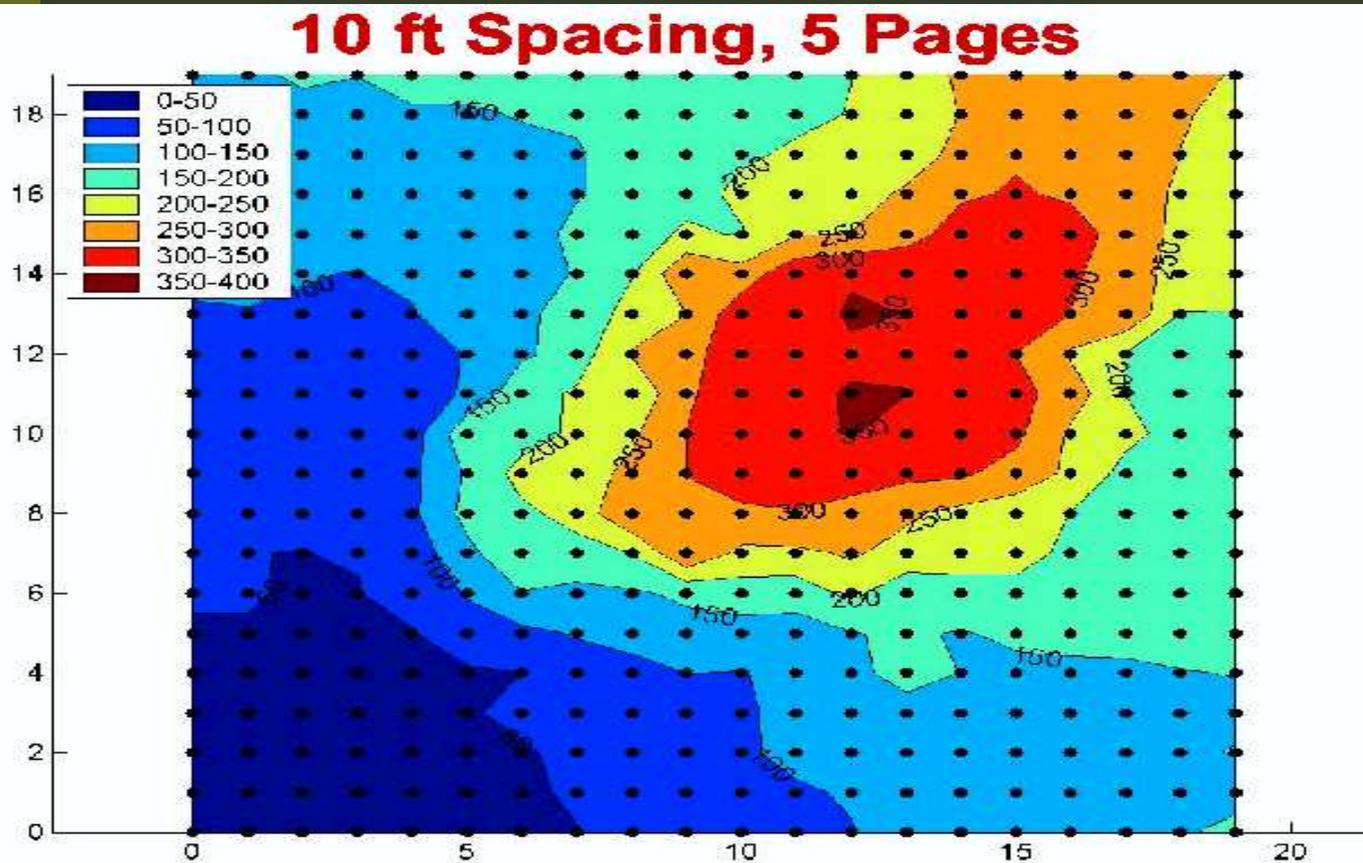
# Simulated Propagation Time for a Sparse Network

Constant rate, wavefront pattern.



# Simulated Propagation Time for a Dense Network

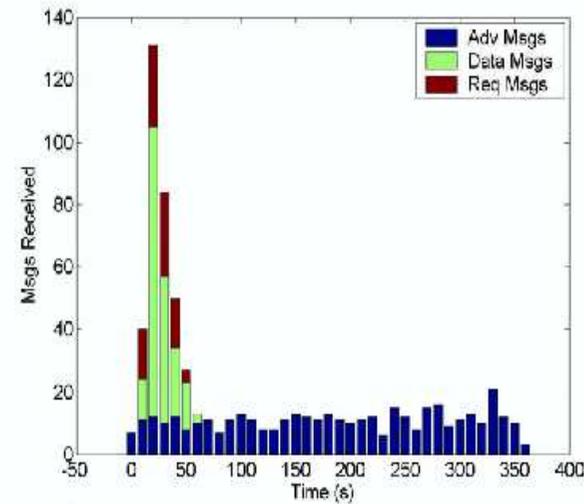
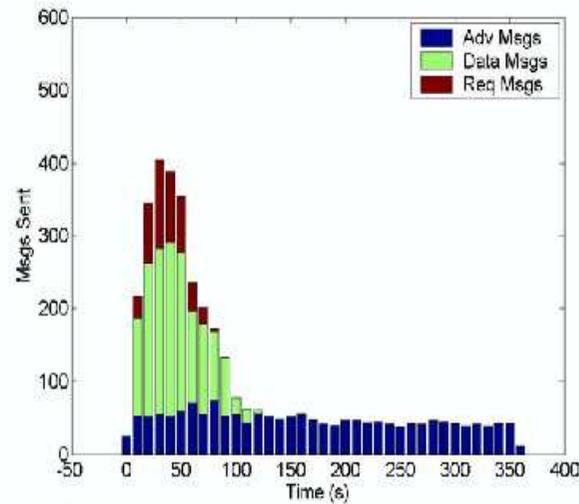
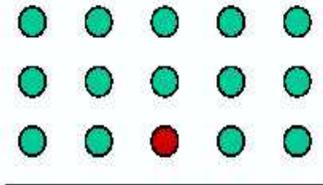
Propagation fast on the edges, slow in the middle.



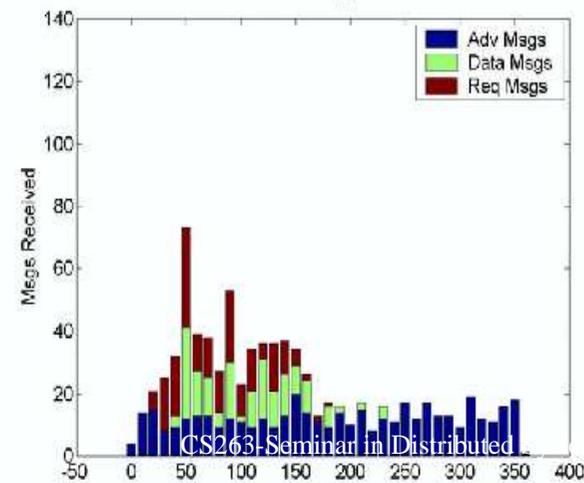
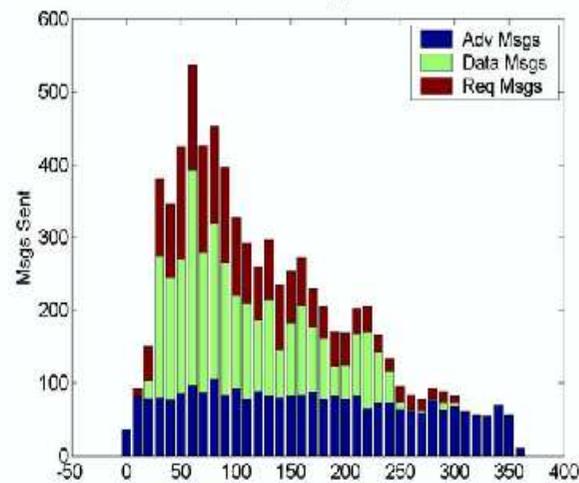
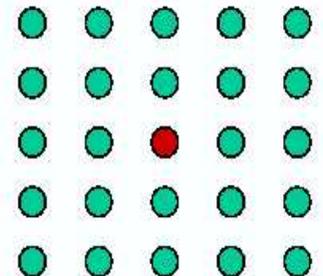
# The Culprit is the Hidden Terminal Problem

Nodes in the center have more neighbors and are more prone to collisions.

## Node Near Edge



## Node In Center



# Lower Bound on Dissemination Rate

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- Dissemination is slower than routing a single message across the network.
- Spatial multiplexing provides great improvement.
- Empirically data redundancy ratio is usually less than 5.

# Conclusions

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- Deluge, a reliable data dissemination protocol for propagating large data objects from a few nodes to many nodes of a multi-hop WSN.
- Epidemic protocol, achieves 90 bytes/second dissemination rate.
- Part of TinyOS 1.1.8.

# Mobile P2P vs WSNs

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Mobile peer-to-peer networks differ from WSNs in a number of ways:

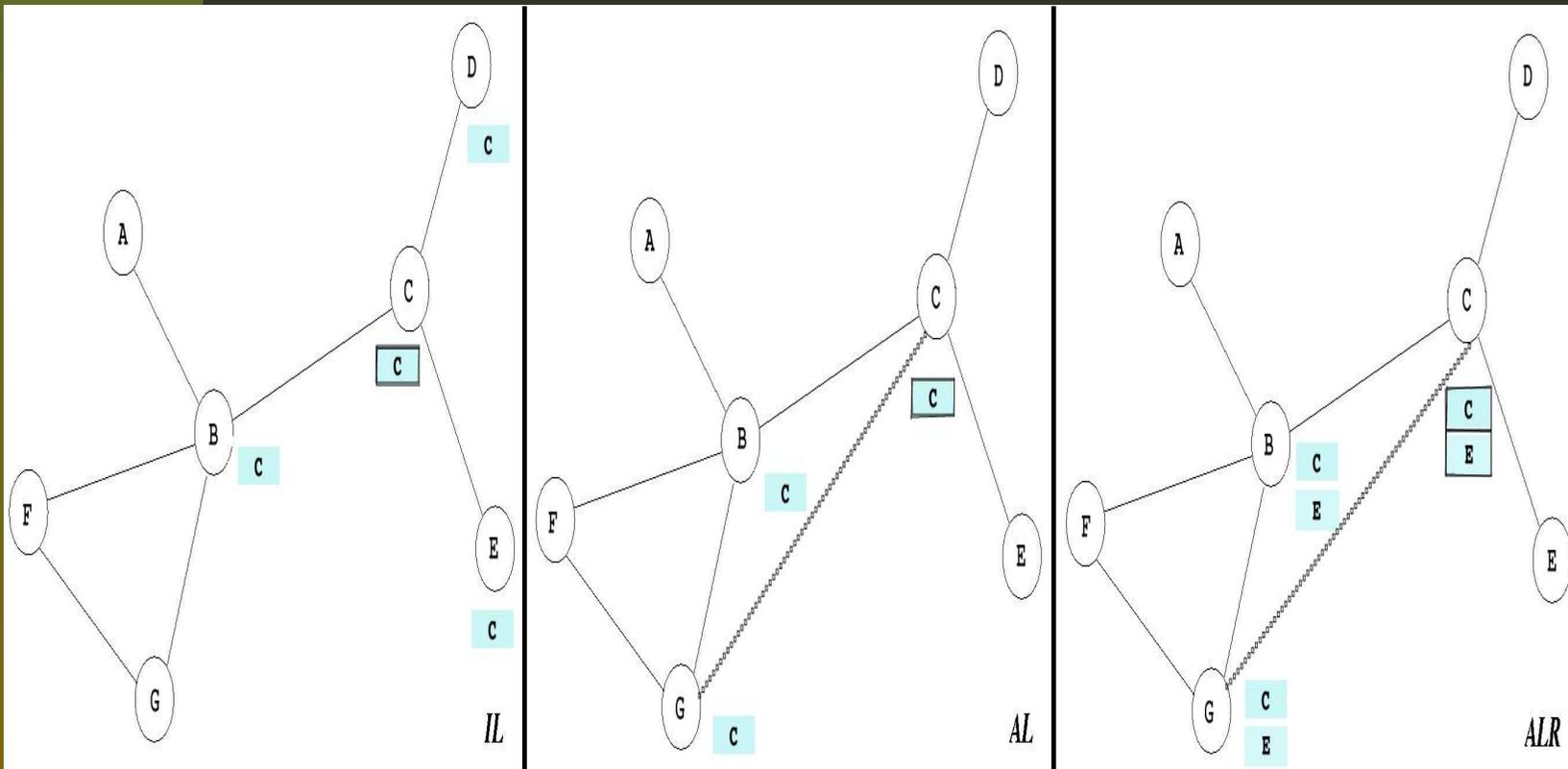
- Built for different purposes.
- Nodes are more powerful, allowing for more complicated protocols.
- Direct connections between the nodes are possible.

# Data Dissemination in Mobile P2P

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- Nodes build content synopses of their data based on Bloom Filters.
- Adaptively disseminate them to the most appropriate nodes of the network.
- Based on the synopses they forward queries to the nodes with the highest probability of providing the desired results.

# Dissemination Strategies



# Thank you!

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Questions/comments?

References:

- The Dynamic Behavior of a Data Dissemination Protocol for Network Programming at Scale, Sensys 2004
- Trickle: A Self-Regulating Algorithm for Code Propagation and Maintenance in Wireless Sensor Networks, NSDI 2004
- Data Dissemination in Mobile Peer-to-Peer Networks, MDM 2005